SDCA without Duality, Regularization, and Individual Convexity

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Minimizing Average-of-Functions

Question: What is the runtime to find w s.t. $F(w) \leq F(w^*) + \epsilon$ where

$$F(w) := \frac{1}{n} \sum_{i=1}^{n} f_i(w)$$

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Classic: Gradient Descent (GD):

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Modern: Stochastic Dual Coordinate Ascent (SDCA):

- Assume: $f_i(w) = \phi_i(w) + \frac{\lambda}{2} \|w\|^2$ and ϕ_i is convex and L smooth
- Runtime: $d \cdot \left(n + \frac{L}{\lambda}\right) \cdot \log\left(\frac{1}{\epsilon}\right)$

Outline

- SDCA without Duality
 - SDCA \in SGD family
 - SGD with a stochastic oracle must be slow
 - SDCA reduces the variance using a stronger oracle
 - A simple convergence proof
- Relaxing the Assumptions
 - Without Explicit Regularization
 - Dependence on Average Smoothness
 - Without Individual Convexity

SDCA without Duality

• Objective:

$$F(w) = \frac{\lambda}{2} ||w||^2 + \frac{1}{n} \sum_{i=1}^{n} \phi_i(w)$$

• At w^* we have $\nabla F(w^*) = 0$:

$$w^* = \frac{1}{\lambda n} \sum_{i=1}^{n} \underbrace{(-\nabla \phi_i(w^*))}_{:=\alpha_i^*}$$

- ullet Primal variable: w
- Pseudo-Dual variables: $\alpha_1, \ldots, \alpha_n$
- Goal: $w^{(t)} \to w^*$ and for every i, $\alpha_i^{(t)} \to \alpha_i^*$

SDCA without Duality

- Initialize: $w^{(0)} = \frac{1}{\lambda n} \sum_{i=1}^{n} \alpha_i^{(0)}$
- For: t = 1, 2, ...
 - Pick $i \in [n]$ at random
 - $\qquad \qquad \text{Primal update: } w^{(t)} = w^{(t-1)} \eta \left(\nabla \phi_i(w^{(t-1)}) + \alpha_i^{(t-1)} \right)$
 - Dual update: $\alpha_i^{(t)} = (1-\beta)\alpha_i^{(t-1)} + \beta\left(-\nabla\phi_i(w^{(t-1)})\right)$ (where $\beta=\eta\lambda n$)

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- Claim: SDCA is an instance of SGD
- Proof:
 - By induction, $\lambda w^{(t-1)} = \frac{1}{n} \sum_{i=1}^n \alpha_i^{(t-1)} = \mathbb{E}_i \alpha_i^{(t-1)}$
 - Therefore:

$$\nabla F(w^{(t-1)}) = \lambda \, w^{(t-1)} + \mathbb{E}_i \, \nabla \phi_i(w^{(t-1)}) = \mathbb{E}_i[\alpha_i^{(t-1)} + \nabla \phi_i(w^{(t-1)})]$$

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SGD with a stochastic oracle must be slow

Theorem

Any algorithm for minimizing F that only accesses the objective using oracle that returns a gradient of a random function and has $\log(1/\epsilon)$ rate must perform $\tilde{\Omega}(n^2)$ iterations

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Proof idea:

• Consider two objectives (in both, $\lambda = 1$): for $i \in \{\pm 1\}$

$$F_i(w) = \frac{1}{2n} \left((n-1) \frac{(w-i)^2}{2} + (n+1) \frac{(w+i)^2}{2} \right)$$

- \bullet A stochastic gradient oracle returns $w\pm i$ w.p. $\frac{1}{2}\pm\frac{1}{2n}$
- Easy to see that $w_i^* = -i/n$, $F_i(0) = 1/2$, $F_i(w_i^*) = 1/2 1/(2n^2)$
- \bullet Therefore, solving to accuracy $\epsilon < 1/(2n^2)$ amounts to determining the bias of the coin

Can we improve SGD ?

A stronger oracle:

- The negative result assumes we only see a gradient of a randomly chosen example
- SDCA relies on a slightly stronger oracle: we also see the index of the chosen example
- This suffices to obtain a significantly faster algorithm
- Main idea: variance reduction

 \bullet SGD update rule: $w^{(t)} = w^{(t-1)} - \eta v$ where $\mathbb{E}[v] = \nabla F(w^{(t-1)})$

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- For SDCA: $v = \nabla \phi_i(w^{(t-1)}) + \alpha_i^{(t-1)}$
- What is the variance?

$$\mathbb{E}[\|v\|^{2}] = \mathbb{E}[\|\nabla\phi_{i}(w^{(t-1)}) - \nabla\phi_{i}(w^{*}) + \nabla\phi_{i}(w^{*}) + \alpha_{i}^{(t-1)}\|^{2}]$$

$$= \mathbb{E}[\|\nabla\phi_{i}(w^{(t-1)}) - \nabla\phi_{i}(w^{*}) + \alpha_{i}^{(t-1)} - \alpha_{i}^{*}\|^{2}]$$

$$\leq 2 \mathbb{E}[\|\nabla\phi_{i}(w^{(t-1)}) - \nabla\phi_{i}(w^{*})\|^{2}] + 2 \mathbb{E}[\|\alpha_{i}^{(t-1)} - \alpha_{i}^{*}\|^{2}]$$

$$\leq 2L \mathbb{E}[\|w^{(t-1)} - w^{*}\|^{2}] + 2 \mathbb{E}[\|\alpha_{i}^{(t-1)} - \alpha_{i}^{*}\|^{2}]$$

• Potential: $C_t = \frac{1}{2L}A_t + \frac{\lambda}{2}B_t$ with $A_t = \mathbb{E}_j \|\alpha_j^{(t)} - \alpha_j^*\|^2$, $B_t = \|w^{(t)} - w^*\|^2$

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- Algebraic Manipulations:

$$\mathbb{E} A_t - (1 - \eta \lambda) A_{t-1} = \eta \lambda \mathbb{E} \left(\| \nabla \phi_i(w^{(t-1)}) - \nabla \phi_i(w^*) \|^2 - (1 - \beta) \|v\|^2 \right)$$

$$\mathbb{E} B_t - B_{t-1} = -2\eta (w^{(t-1)} - w^*)^\top \nabla F(w^{(t-1)}) + \eta^2 \mathbb{E} \|v\|^2$$

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- Summing with weights $(\frac{1}{2L}, \frac{\lambda}{2})$ cancels the $\mathbb{E} \|v\|^2$ term and gives

$$C_t - (1 - \eta \lambda) C_{t-1} \le \eta \lambda \left(\frac{1}{2L} \mathbb{E} \| \nabla \phi_i(w^{(t-1)}) - \nabla \phi_i(w^*) \|^2 - \epsilon_{t-1} \right)$$

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 $\bullet \ \phi_i \text{ is L-smooth and convex} \Rightarrow \ \mathbb{E}_i \, \|\nabla \phi_i(w^{(t-1)}) - \nabla \phi_i(w^*)\|^2 \leq 2L\epsilon_{t-1}$

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- Therefore: $\mathbb{E} C_t \leq (1 \eta \lambda) \mathbb{E} C_{t-1} \leq (1 \eta \lambda)^t C_0$

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GD vs SDCA

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SDCA Without Explicit Regularization

Original objective:

$$F(w) = \frac{1}{n} \sum_{i=1}^{n} \phi_i(w)$$

Rewrite the objective as

$$F(w) = \frac{1}{n+1} \sum_{i=1}^{n+1} \phi_i(w) + \frac{\lambda}{2} ||w||^2$$

where

- For $i \leq n$, $\phi_i(w) = \frac{n+1}{n} f_i(w)$
- $\phi_{n+1}(w) = \frac{-\lambda(n+1)}{2} ||w||^2$

Dependence on Average Smoothness

- Assume that ϕ_i is L_i -smooth
- Let $\bar{L} = \mathbb{E}_i L_i$
- Sample $i \sim q$ where $q_i = \frac{L_i + \bar{L}}{2n\bar{L}}$
- ullet Convergence rate now depends on $ar{L}$ instead of on $\max_i L_i$

SDCA without Individual Convexity

- Remove the assumption that ϕ_i is convex and only assume that F is λ -strongly convex
- The bound $\mathbb{E}_i \, \| \nabla \phi_i(w^{(t-1)}) \nabla \phi_i(w^*) \|^2 \leq 2L\epsilon_{t-1}$ no longer holds
- Instead, $\mathbb{E}_i \| \nabla \phi_i(w^{(t-1)}) \nabla \phi_i(w^*) \|^2 \le L^2 \| (w^{(t-1)} w^*) \|^2$
- ullet Yields a runtime of $d\cdot \left(n+\left(rac{L}{\lambda}
 ight)^2
 ight)\cdot \log\left(rac{1}{\epsilon}
 ight)$
- \bullet Using acceleration gives runtime of $\tilde{O}\left(d\cdot(n+n^{3/4}\sqrt{\bar{L}/\lambda})\right)$
- ullet Compare to the convex case: $\tilde{O}\left(d\cdot(n+n^{1/2}\sqrt{\bar{L}/\lambda})
 ight)$

Summary

- SDCA without duality as a variance reduced SGD
- Simpler proof
- Relaxing the assumptions on individual functions
- ullet Open: is the extra $n^{1/4}$ factor necessary ?