## Algorithms - Exercise 5

Due Wednesday 1/12 24:00

- 1. (a) Show that  $1+y \leq e^y$  for any real y with strict inequality for  $y \neq 0$ 
  - (b) Show that  $1 \epsilon \le (1 \frac{\epsilon}{n})^n$  for  $n \ge 1$  and  $0 < \epsilon \le 1$
  - (c) Show that for any three sets A, B, C one has

$$A\cap (B\setminus C)=(A\setminus C)\cap (B\setminus C)$$

- 2. Show that the  $2 \frac{1}{k}$  bound on the approximation ratio of the algorithm for assignment distribution problem shown in class is tight. That is, show for any k there exists an instance of the problem on which the algorithm gives exactly  $2 \frac{1}{k}$ -approximate solution.
- 3. How would you modify the approximation scheme for Subset-Sum presented in class to find a good approximation to the smallest value not less than t that is a sum of some subset of the given input list?
- 4. Show that the usual approximation algorithm for Vertex-Cover doesn't work in a graph with weighted vertexes, when the problem is to find a cover with the minimal weight.
- 5. Suppose we generalize the set-covering problem so that each set  $S_i$  in the family F has an associated weight  $w_i$  and the weight of a cover C is  $\sum_{S_i \in C} w_i$ . We wish to determine a minimum-weight cover. Change the algorithm given in class for Set-Cover and show that the bound on approximation ratio still holds.