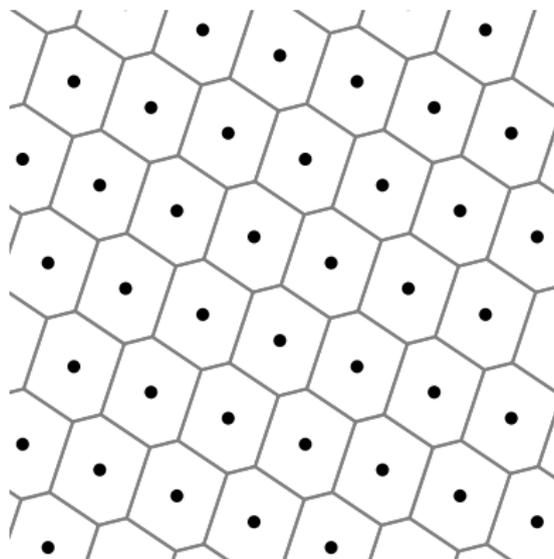


# A Simple Proof for the Existence of “Good” Pairs of Nested Lattices

Or Ordentlich  
Joint work with Uri Erez

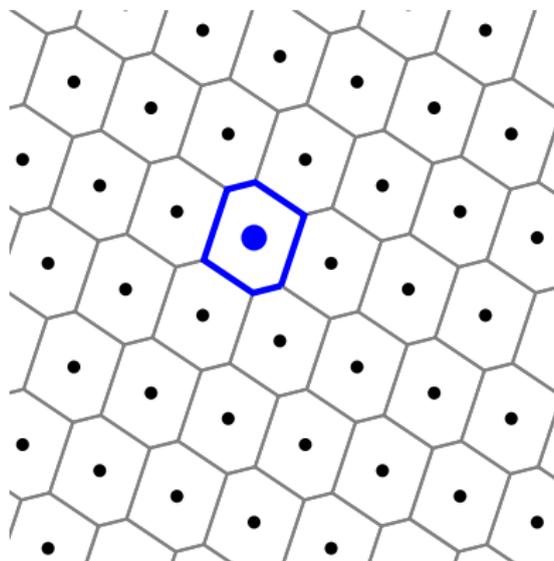
November 15th, IEEEI 2012  
Eilat, Israel

# What are lattices?



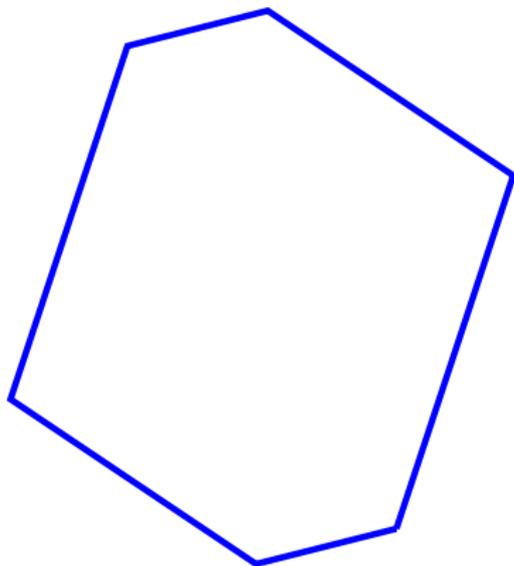
- A lattice  $\Lambda$  is a discrete subgroup of  $\mathbb{R}^n$  closed under addition and reflection.

# Lattice definitions



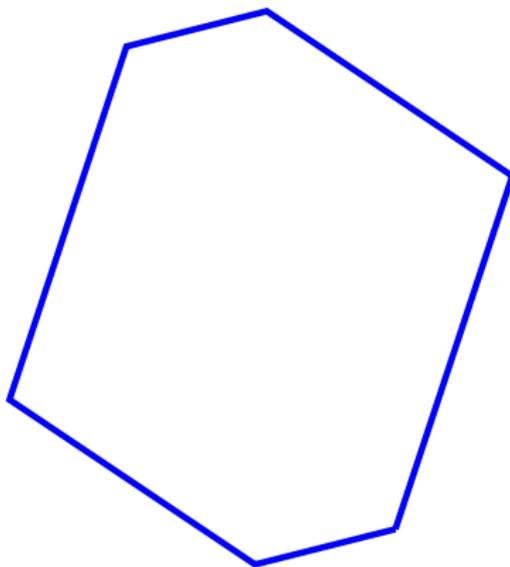
- The Voronoi region  $\mathcal{V}$  of a lattice point is the set of all points in  $\mathbb{R}^n$  which are closest to it.

# Lattice definitions



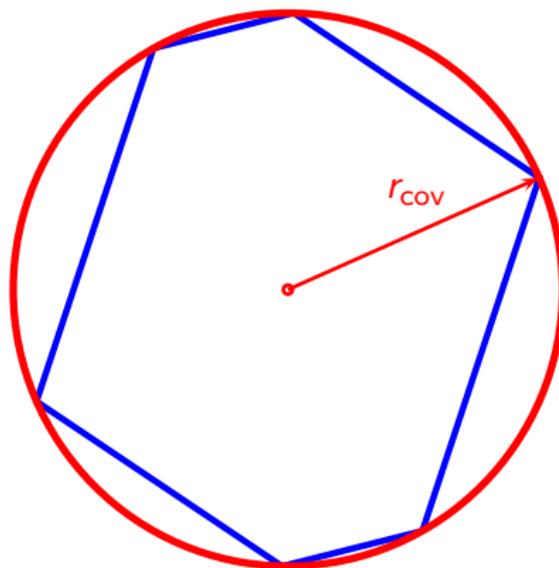
- The Voronoi region  $\mathcal{V}$  of a lattice point is the set of all points in  $\mathbb{R}^n$  which are closest to it.

# Lattice definitions



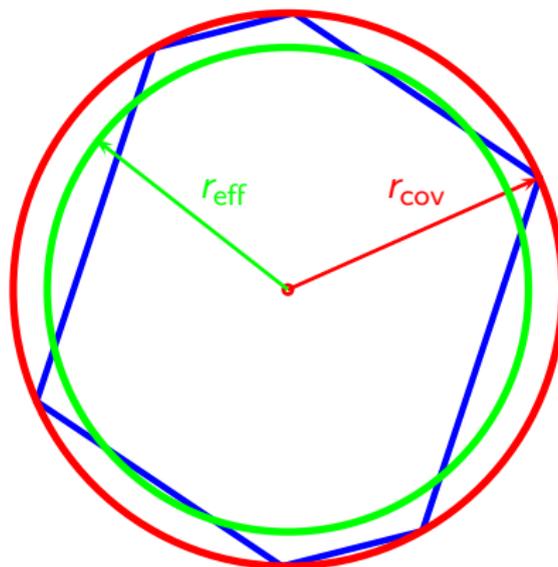
- A sequence of lattices is good for AWGN coding if the probability that an AWGN (with appropriate variance) is not contained in  $\mathcal{V}$  vanishes with the dimension.

# Lattice definitions



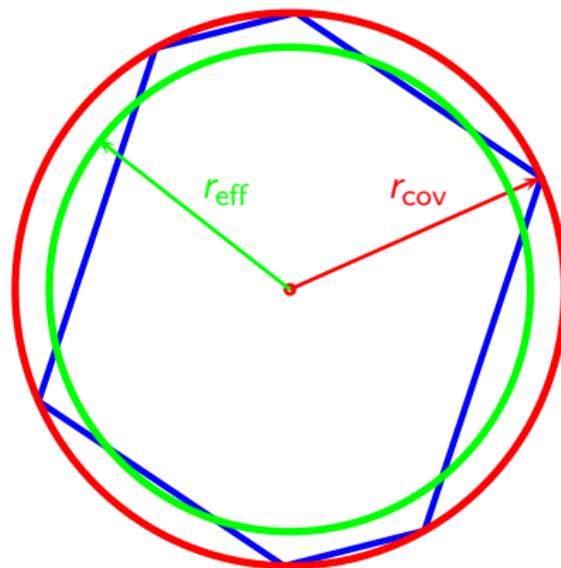
- The covering radius is the radius of the smallest ball that contains the Voronoi region  $\mathcal{V}$ .

# Lattice definitions



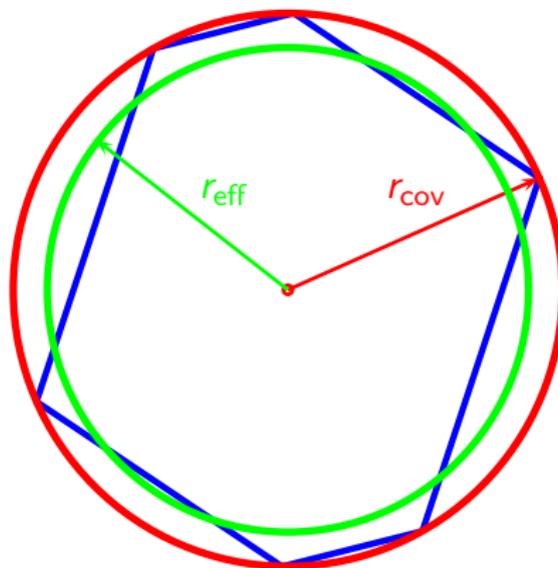
- The effective radius is the radius of a ball with the same volume as the Voronoi region  $\mathcal{V}$ .

# Lattice definitions



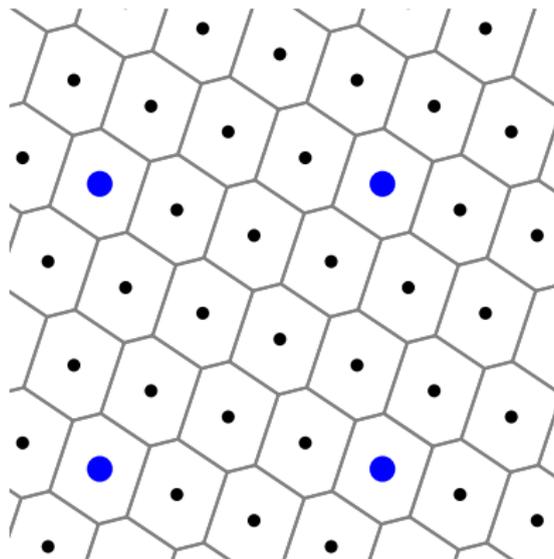
- A sequence of lattices is good for covering (“Rogers-good”) if  $\frac{r_{\text{eff}}}{r_{\text{cov}}} \rightarrow 1$  as the lattice dimension grows.

# Lattice definitions



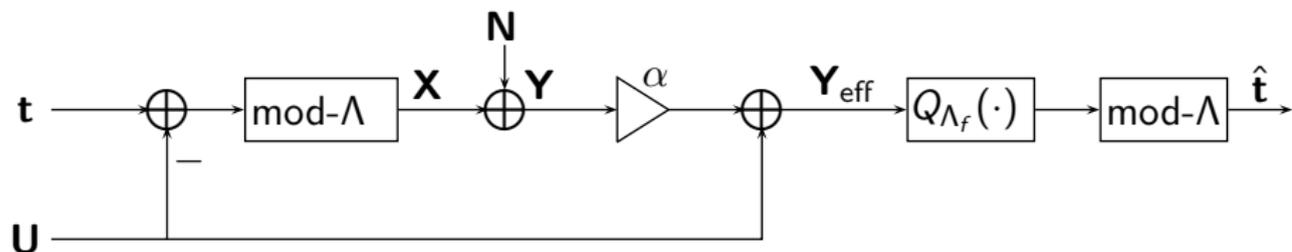
- A sequence of lattices is good for MSE quantization if for a random dither  $\mathbf{U}$  uniformly distributed over  $\mathcal{V}$  we have  $\mathbb{E}\|\mathbf{U}\|^2 \rightarrow r_{\text{eff}}^2$  as the lattice dimension grows.

# Lattice definitions



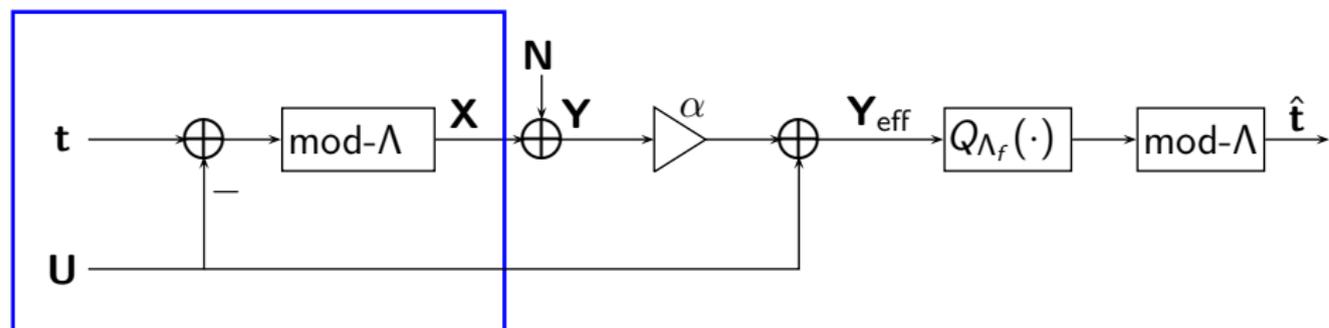
- A lattice  $\Lambda_c$  is nested in  $\Lambda$  if  $\Lambda_c \subset \Lambda$ .

# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]



- The scheme uses a pair of nested lattice  $\Lambda \subset \Lambda_f$  and a dither  $U$ .

# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]



Tx:

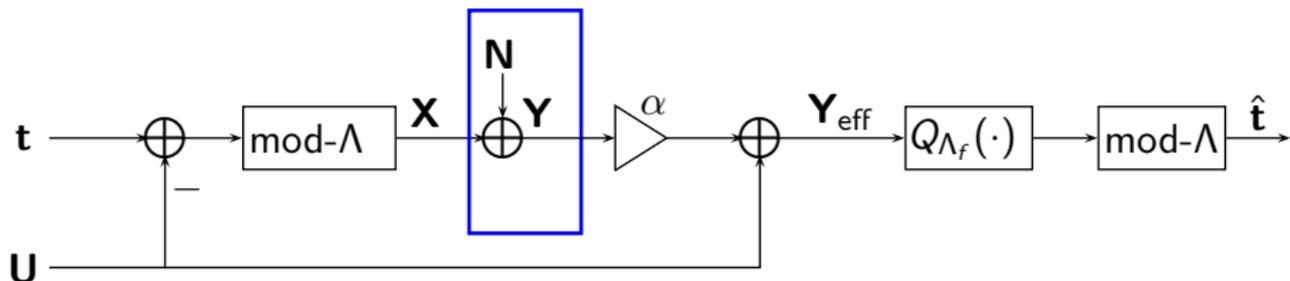
- The transmitted signal is

$$\mathbf{X} = [\mathbf{t} - \mathbf{U}] \bmod \Lambda$$

- $\mathbf{X}$  is uniformly distributed over the Voronoi region of  $\Lambda$ .
- The average transmission power is the second moment of  $\Lambda$ .

$\Lambda$  needs to be good for MSE quantization

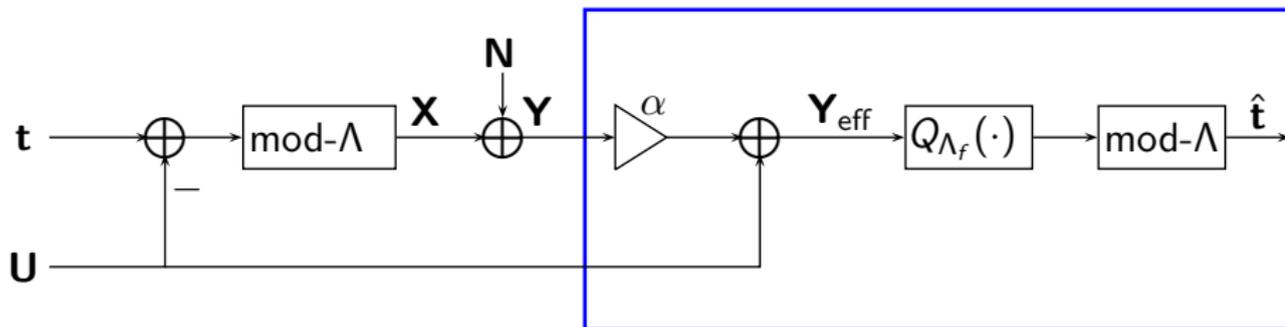
# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]



AWGN Channel:

$$Y = X + N$$

# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]



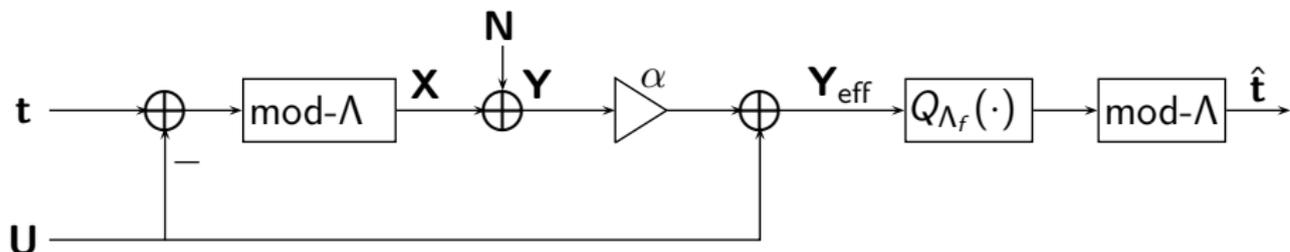
Rx:

$$\begin{aligned} \mathbf{Y}_{\text{eff}} &= \alpha \mathbf{Y} + \mathbf{U} \\ &= \mathbf{X} + \mathbf{U} + (\alpha - 1)\mathbf{X} + \alpha \mathbf{N} \\ &= \mathbf{t} - \mathbf{U} + \lambda + \mathbf{U} + (\alpha - 1)\mathbf{X} + \alpha \mathbf{N} \\ &= \mathbf{t} + \lambda + \mathbf{Z}_{\text{eff}}, \end{aligned}$$

where  $\lambda \in \Lambda$  and  $\mathbf{Z}_{\text{eff}} = (\alpha - 1)\mathbf{X} + \alpha \mathbf{N}$ .

Decoding is correct if  $\mathbf{Z}_{\text{eff}} \in \mathcal{V}_f$

# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]

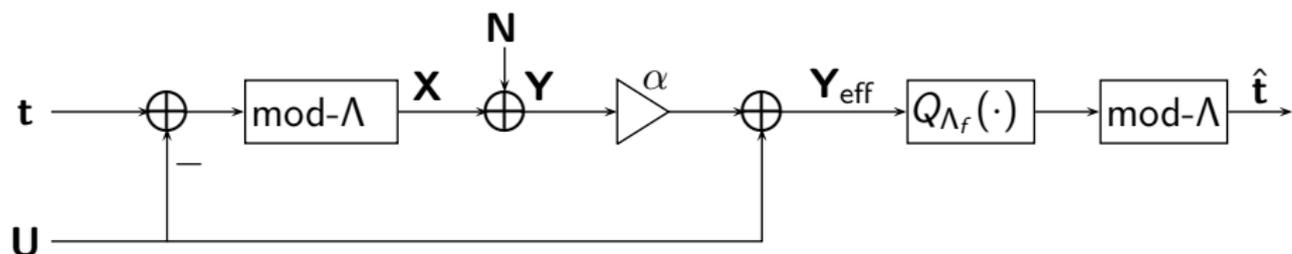


To approach capacity with nearest neighbor decoding we need:

- The coarse lattice  $\Lambda$  to be good for MSE quantization.
- The fine lattice  $\Lambda_f$  to be good for coding (with NN decoding) in the presence of effective noise  $\mathbf{Z}_{\text{eff}} = (\alpha - 1)\mathbf{X} + \alpha\mathbf{N}$ .

Note that  $\mathbf{Z}_{\text{eff}}$  depends on the coarse lattice.

# The mod- $\Lambda$ transmission scheme [Erez-Zamir IT04]



To be more formal, we want:

- $\frac{1}{n}\mathbb{E}\|\mathbf{U}\|^2 \rightarrow \frac{1}{2\pi e}\text{Vol}(\Lambda)^{2/n}$  (or alternatively  $G(\Lambda) \rightarrow \frac{1}{2\pi e}$ ).
- If  $V(\Lambda_f)^{\frac{2}{n}} > 2\pi e \frac{1}{n}\mathbb{E}\|\mathbf{Z}_{\text{eff}}\|^2$ ,

$$\Pr ([Q_{\Lambda_f}(\mathbf{t} + \mathbf{Z}_{\text{eff}})] \bmod \Lambda \neq \mathbf{t} \bmod \Lambda) \rightarrow 0.$$

If the two conditions are satisfied the scheme achieves any rate below

$$R = \frac{1}{n} \log \left( \frac{\text{Vol}(\Lambda)}{\text{Vol}(\Lambda_f)} \right) = \frac{1}{2} \log \left( \frac{\text{SNR}}{\frac{1}{n}\mathbb{E}\|\mathbf{Z}_{\text{eff}}\|^2} \right)$$

# How to construct “good” pairs of nested lattices?

	Binary linear code	Lattice
Single		
Nested		

# How to construct “good” pairs of nested lattices?

	Binary linear code	Lattice
Single	Draw $G \in \mathbb{Z}_2^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$	
Nested		

# How to construct “good” pairs of nested lattices?

	Binary linear code	Lattice
Single	Draw $G \in \mathbb{Z}_2^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$	Draw $G \in \mathbb{Z}_p^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_p$ . Construct a linear code $\mathcal{C}$ with this generating matrix. Set $\Lambda = p^{-1}\mathcal{C} + \mathbb{Z}^n$
Nested		

# How to construct “good” pairs of nested lattices?

	Binary linear code	Lattice
Single	Draw $G \in \mathbb{Z}_2^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$	Draw $G \in \mathbb{Z}_p^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_p$ . Construct a linear code $\mathcal{C}$ with this generating matrix. Set $\Lambda = p^{-1}\mathcal{C} + \mathbb{Z}^n$
Nested	Draw $G_f = \begin{bmatrix} \mathbf{G} \\ - - - \\ \mathbf{G}' \end{bmatrix}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$ . $G$ generates the coarse code and $G_f$ the fine code	

# How to construct “good” pairs of nested lattices?

	Binary linear code	Lattice
Single	Draw $G \in \mathbb{Z}_2^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$	Draw $G \in \mathbb{Z}_p^{K \times N}$ with entries i.i.d. and uniform over $\mathbb{Z}_p$ . Construct a linear code $\mathcal{C}$ with this generating matrix. Set $\Lambda = p^{-1}\mathcal{C} + \mathbb{Z}^n$
Nested	Draw $G_f = \begin{bmatrix} \mathbf{G} \\ - - - \\ \mathbf{G}' \end{bmatrix}$ with entries i.i.d. and uniform over $\mathbb{Z}_2$ . $G$ generates the coarse code and $G_f$ the fine code	?

# How to construct “good” pairs of nested lattices?

## Previous works

- Find a lattice  $\Lambda$  which is good for covering, and has a generating matrix  $F$ .
- Draw  $G \in \mathbb{Z}_p^{K \times N}$  with entries i.i.d. and uniform over  $\mathbb{Z}_p$ . Construct a linear code  $\mathcal{C}$  with this generating matrix.
- Set  $\Lambda_f = F \cdot (p^{-1}\mathcal{C} + \mathbb{Z}^n) = p^{-1}FC + F\mathbb{Z}^n$ .

## This work

- Draw  $G_f = \begin{bmatrix} \mathbf{G} \\ - - - \\ \mathbf{G}' \end{bmatrix} \in \mathbb{Z}_p$  with entries i.i.d. and uniform over  $\mathbb{Z}_p$ .
- Construct a linear code  $\mathcal{C}$  from the generating matrix  $G$ , and a linear code  $\mathcal{C}_f$  from the generating matrix  $G_f$ .
- Set  $\Lambda = p^{-1}\mathcal{C} + \mathbb{Z}^n$ ,  $\Lambda_f = p^{-1}\mathcal{C}_f + \mathbb{Z}^n$ .

# Differences from previous work

## Current approach vs. previous work

- A simpler ensemble to analyze.
- A basic proof that makes no use of previous results from geometry of numbers.
- The coarse lattice only has to be good for MSE quantization (not necessarily for covering).

# Differences from previous work

## Current approach vs. previous work

- A simpler ensemble to analyze.
- A basic proof that makes no use of previous results from geometry of numbers.
- The coarse lattice only has to be good for MSE quantization (not necessarily for covering).

## Historical note:

- When Erez and Zamir's work was published it was known that good covering lattices exist, but it was not known that Construction A lattices are usually good for covering.  
⇒ The two-step construction was needed.
- Erez and Zamir studied the error exponents the mod- $\Lambda$  scheme achieves. For that purpose it is important that the coarse lattice will be good for covering. For capacity, goodness for quantization suffices.

# Proof outline

- We show that w.h.p. the coarse lattice  $\Lambda$  is good for MSE quantization.
  - We show that w.h.p. the fine lattice  $\Lambda_f$  is good for coding in the presence of noise that rarely leaves a ball.
- ⇒ Most members of the ensemble are “good” pairs of nested lattices.
- We show that for a coarse lattice  $\Lambda$  which is good for MSE quantization the effective noise

$$\mathbf{Z}_{\text{eff}} = (\alpha - 1)\mathbf{X} + \alpha\mathbf{N}$$

rarely leaves a ball.

# Proof outline - Goodness for MSE quantization

- Our ensemble induces a distribution on the lattice points that is “almost” i.i.d. and “uniform” on  $\mathbb{R}^n$  with point density  $1/(V_n r_{\text{eff}}^n)$ .
- For any  $\mathbf{x} \in \mathbb{R}^n$ :  $|\Lambda \cap \mathcal{B}(\mathbf{x}, \sqrt{nD})| \sim \left(\frac{\sqrt{nD}}{r_{\text{eff}}}\right)^n$ .
  - $\implies$  For  $\sqrt{nD} > r_{\text{eff}}$  almost surely  $|\Lambda \cap \mathcal{B}(\mathbf{x}, \sqrt{nD})| > 1$ .
  - $\implies$  If  $\sqrt{nD} > r_{\text{eff}}$ , almost surely a point  $\mathbf{x} \in \mathbb{R}^n$  is covered by a lattice point with distance  $< \sqrt{nD}$ .
- Since  $\Lambda$  is nested within the cubic lattice, the covering radius is bounded.
  - $\implies$  For  $\sqrt{nD} > r_{\text{eff}}$  the average MSE distortion  $\Lambda$  achieves is smaller than  $nD$ .
- Setting  $r_{\text{eff}}$  slightly smaller than  $\sqrt{nD}$ , a dither  $\mathbf{U}$  uniformly distributed over  $\mathcal{V}$  satisfies

$$\frac{1}{n} \mathbb{E} \|\mathbf{U}\|^2 \leq D \approx \frac{r_{\text{eff}}^2}{n} = \frac{1}{n V_n^{2/n}} V(\Lambda)^{2/n} \approx \frac{V(\Lambda)^{2/n}}{2\pi e}$$

# Proof outline - effective noise rarely leaves a ball

## Lemma:

Let  $\mathbf{U}$  be a dither from a lattice  $\Lambda$  which is good for MSE quantization and has effective radius  $r_{\text{eff}}$ . Then, for any  $\epsilon > 0$  and  $n$  large enough

$$\Pr(\mathbf{U} \notin \mathcal{B}(\mathbf{0}, (1 + \epsilon)r_{\text{eff}})) < \epsilon$$

- The pdf of a random vector uniform over a ball is upper bounded by that of an AWGN. [Erez-Zamir 04]
- Combining with our lemma, w.p.  $1 - \epsilon$  we have that  $\mathbf{U}$  is approximately Gaussian .

With probability  $1 - \epsilon$  the effective noise  $\mathbf{Z}_{\text{eff}} = (1 - \alpha)\mathbf{U} + \alpha\mathbf{N}$  is Gaussian.  
 $\implies$  It suffices that the fine lattice will be good for AWGN channels.

Showing that Construction A lattices are good for AWGN channels is straightforward.

# Corollaries and extensions

## Chains of nested lattices

A similar construction results in a chain of nested lattice codes which are all good for MSE quantization and AWGN coding.

## Cubic shaping lattice

If the coarse lattice is cubic (no shaping) and the fine lattice is good for AWGN coding, the mod- $\Lambda$  scheme can achieve any rate satisfying

$$R < \frac{1}{2} \log(1 + \text{SNR}) - \frac{1}{2} \log\left(\frac{2\pi e}{12}\right).$$

## Additive non-Gaussian noise

For any additive i.i.d. noise channel the mod- $\Lambda$  scheme can achieve any rate satisfying  $R < \frac{1}{2} \log(1 + \text{SNR})$  with nearest-neighbor decoding (followed by mod- $\Lambda$ ).

This is reminiscent of Lapidoth 96.